



cmp - solution

There are several solutions for this problem. The best solution the Scientific Committee has uses $T = 10$, but we believe that $T = 9$ is also achievable and deserves more than 100 points.

- *A simple $T = 17$ solution:*

We consider a binary tree that has 4096 leaves.

Remember(a): we set all the nodes from leaf “a” to the root to 1.

maxA = 12

Compare(b): Our goal is to determine the lowest common ancestor (LCA) for leaf a and leaf b. Since the whole path from the root to the leaf is set to 1 we can do a binary search. Once we have the LCA we can return the -1, 0 or 1.

maxB = 5

The next solutions use **one-hot encoding**.

This encoding uses N bits for representing a number from 0 to $N-1$. For a given number X we only set the bit X to 1 and leave the rest set to 0. The advantage of this method is that we only need to write 1 bit. While reading a value could take $N - 1$ reads (we look for the bit set to 1), comparing is a little faster. It takes $(N-1)/2$ reads. We only need to search from the current position to the closest end (0 or $N-1$).

- *Some tree solutions ($T = 13$):*

We change the simple binary tree from the previous solution. We still have at least 4096 leaves, but the internal nodes will have a variable number of children. For instance let's assume that we have 6 levels and the number of children for each level is A, B, C, D, E and F. We chose the degree for each level such that $A*B*C*D*E*F \geq 4096$.

To encode a value in a node we use the **one-hot** encoding.



Let's consider $A = B = C = D = E = F = 4$

Remember(a): we set all the nodes from leaf “a” to the root to 1.

$\max A = 6$, since we have 6 nodes and encoding takes 1 `bit_set()` call.

Compare(b): Considering the same representation for b again we look for LCA. However this time we do a top-down traversal. As soon as the bit we expect to be set to 1 is 0 we try to determine if b is higher or lower.

Worst case is when the last bit is not set 1. In this case we need to read an additional bit.

$\max B = 7$, since we have 6 nodes and an additional `bit_get()`.

- *Further optimizations ($T = 12$):*

For the previous solution we observe that the worst case is achieved when we miss the bit in the last node.

If we choose $A = 7, B = 6, C = 5, D = 4, E = 3, F = 2$ then $\max B$ becomes 6. This solution also works if comparing a one-hot encoding is done in $N - 1$ reads and not $(N-1)/2$.

- *Breaking the $\max A - \max B$ symmetry ($T = 10$):*

Instead of 6 levels for our tree we chose 4 levels of degree: 12, 10, 8, 6.

$\max A$ becomes 4 and $\max B$ becomes 6.

This solution uses the $T = 12$ optimization.

For a given set of maxim branching factors the formula for T is:

$num_levels + \max(branching_factor_on_level[i]/2 + 1 + i \text{ for } i \text{ in range}(0, num_levels))$

with $product(branching_factor_on_level(i)) > 4095$



- *Mixed-radix*

All the above tree solutions don't really need a tree. For instance let's pick the $T = 10$ above solution.

We can encode the value from the root in the bits from 1 to 12, the value from the node on the second level in the memory from 13 to 23 and so on.