



## cmp - solution

There are several solutions for this problem. The best solution the Scientific Committee has uses  $T = 10$ , but we believe that  $T = 9$  is also achievable and deserves more than 100 points.

- *A simple  $T = 17$  solution:*

We consider a binary tree that has 4096 leaves.

Remember(a): we set all the nodes from leaf “a” to the root to 1.

maxA = 12

Compare(b): Our goal is to determine the lowest common ancestor (LCA) for leaf a and leaf b. Since the whole path from the root to the leaf is set to 1 we can do a binary search. Once we have the LCA we can return the -1, 0 or 1.

maxB = 5

The next solutions use **one-hot encoding**.

This encoding uses  $N$  bits for representing a number from 0 to  $N-1$ . For a given number  $X$  we only set the bit  $X$  to 1 and leave the rest set to 0. The advantage of this method is that we only need to write 1 bit. While reading a value could take  $N - 1$  reads (we look for the bit set to 1), comparing is a little faster. It takes  $(N-1)/2$  reads. We only need to search from the current position to the closest end (0 or  $N-1$ ).

- *Some tree solutions ( $T = 13$ ):*

We change the simple binary tree from the previous solution. We still have at least 4096 leaves, but the internal nodes will have a variable number of children. For instance let's assume that we have 6 levels and the number of children for each level is A, B, C, D, E and F. We chose the degree for each level such that  $A*B*C*D*E*F \geq 4096$ .

To encode a value in a node we use the **one-hot** encoding.



Let's consider  $A = B = C = D = E = F = 4$

*Remember(a)*: we set all the nodes from leaf “a” to the root to 1.

$\max A = 6$ , since we have 6 nodes and encoding takes 1 `bit_set()` call.

*Compare(b)*: Considering the same representation for b again we look for LCA. However this time we do a top-down traversal. As soon as the bit we expect to be set to 1 is 0 we try to determine if b is higher or lower.

Worst case is when the last bit is not set 1. In this case we need to read an additional bit.

$\max B = 7$ , since we have 6 nodes and an additional `bit_get()`.

- *Further optimizations ( $T = 12$ ):*

For the previous solution we observe that the worst case is achieved when we miss the bit in the last node.

If we choose  $A = 7, B = 6, C = 5, D = 4, E = 3, F = 2$  then  $\max B$  becomes 6. This solution also works if comparing a one-hot encoding is done in  $N - 1$  reads and not  $(N-1)/2$ .

- *Breaking the  $\max A - \max B$  symmetry ( $T = 10$ ):*

Instead of 6 levels for our tree we chose 4 levels of degree: 12, 10, 8, 6.

$\max A$  becomes 4 and  $\max B$  becomes 6.

This solution uses the  $T = 12$  optimization.

For a given set of maxim branching factors the formula for T is:

$num\_levels + \max(branching\_factor\_on\_level[i]/2 + 1 + i \text{ for } i \text{ in range}(0, num\_levels))$

with  $product(branching\_factor\_on\_level(i)) > 4095$



- *Mixed-radix*

All the above tree solutions don't really need a tree. For instance let's pick the  $T = 10$  above solution.

We can encode the value from the root in the bits from 1 to 12, the value from the node on the second level in the memory from 13 to 23 and so on.